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APPLICATION FOR LETTERS PATENT

**Smart Card with Volatile Memory File Subsystem**

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1 **TECHNICAL FIELD**

2 This invention relates to integrated circuit (IC) modules, such as smart  
3 cards. More particularly, this invention relates to file subsystems used in IC  
4 modules.

5  
6 **BACKGROUND**

7 Integrated circuit (IC) modules, such as smart cards, are widely used in  
8 many different areas. One common use of smart cards is electronic commerce.  
9 Smart cards securely store electronic currency that may be used to purchase  
10 services or products electronically over a network or at a point-of-purchase  
11 terminal (e.g., kiosk, vending machine, phone, toll booth, etc.). Other areas in  
12 which smart cards are deployed include physical and logical access control of  
13 buildings and computer systems, storage of critical information (e.g., medical  
14 records), unscrambling of cable or satellite "premium" signals, and ticketless  
15 travel.

16 Smart cards have evolved technologically to include a processor and  
17 multiple forms of memory, including both volatile memory (e.g., random access  
18 memory (RAM), and nonvolatile memory (e.g., FLASH, read only memory  
19 (ROM), electrically erasable programmable read only memory EEPROM, etc.).  
20 As smart cards grew more sophisticated, they began supporting an operating  
21 system and multiple applications. One exemplary smart card operating system is  
22 Microsoft's Windows® for Smart Card operating system.

23 With the Windows® operating system, smart card applications can declare  
24 "volatile" variables that are stored in RAM. Today, those volatile variables are  
25 bound to the application and are disposed once the application terminates.

1 Currently, for the Windows® for Smart Card operating system, there is no means  
2 to store and retrieve data in volatile memory for access by multiple applications.  
3 Another smart card operating system, called "MULTOS", does allow multiple  
4 applications to store and retrieve data in volatile memory, but its scheme uses a  
5 designated memory area block and associates digital certificates to share that  
6 space across applications in a controlled manner. This scheme is very complex  
7 and not simple to use or implement.

8 Sharing in RAM, however, poses another problem in terms of security.  
9 There is some concern that sensitive data in volatile memory might be accessed by  
10 a rogue application and used in a manner not anticipated by the smart card  
11 designer or user.

12 Accordingly, there is a need for technology that renders data in volatile  
13 memory available to multiple applications in a simple way, but in a secure fashion  
14 to protect against unintentional usage by rogue or malicious applications.

## 16 SUMMARY

17 An integrated circuit (IC) module allows volatile data generated by  
18 applications to be stored within volatile data files in the volatile memory. A file  
19 system tracks the location of all data files as residing in either volatile memory or  
20 nonvolatile memory. The file system then facilitates access to the volatile data  
21 files in volatile memory in a manner similar to accessing nonvolatile data files in  
22 nonvolatile memory.

23 The file system exposes a set of application program interfaces (APIs) to  
24 allow applications to access the data files. The same APIs are used to access both  
25 volatile data files and nonvolatile data files. When an application requests access

1 to a data file, the file system initially determines whether the application is  
2 authorized to gain access to the data file. If it is, the file system next determines  
3 whether the data file resides in volatile memory or nonvolatile memory. Once the  
4 memory region is identified, the file system identifies the physical location of the  
5 data file within that memory region.

6 In one implementation, the file system employs a memory region directory  
7 to identify the memory region in which the data files are located (i.e., RAM,  
8 ROM, EEPROM, etc.). Once the region is identified, one or more separate file  
9 location specifiers specify the physical locations of the data files. As an example,  
10 the file system might include one file allocation table (FAT) to find data file(s) in  
11 the EEPROM and a second FAT to find data file(s) in the RAM. The second FAT  
12 is temporarily constructed and stored in RAM to specify locations of volatile data  
13 files within RAM.

14 In another implementation, the API functions may be modified to specify  
15 the memory region in which a data file is located. By modifying the functions, an  
16 application may pass in a parameter specifying the memory region in which the  
17 data file is currently stored.

18 The file system enables multiple applications to share the same data in  
19 volatile memory by accessing the volatile data files via the file system. By  
20 extending the file system to manage data in the volatile memory, the volatile data  
21 files are protected by the access control mechanisms of the file system so that only  
22 authorized applications can perform file operations on the volatile data files. As a  
23 second layer of protection, the volatile data files and RAM-based FAT are  
24 automatically destroyed when power is removed from the IC module and any  
25 reference to such files kept in the memory region directory is deleted upon

1 initialization of the IC module when power is restored. As a result, any sensitive  
2 data that may have been stored in RAM is never available to rogue applications.

### 3 4 **BRIEF DESCRIPTION OF THE DRAWINGS**

5 Fig. 1 is a block diagram of an integrated circuit (IC) module embodied as a  
6 smart card.

7 Fig. 2 is a logical representation of the file system on the IC module.

8 Fig. 3 is a flow diagram of a process for operating the file system of the IC  
9 module.

10 Fig. 4 is a flow diagram of a process for opening a file that might reside in  
11 volatile or nonvolatile memory.

### 12 13 **DETAILED DESCRIPTION**

14 An integrated circuit (IC) module is designed with a file system that  
15 facilitates access to data files in nonvolatile memory and to volatile data files in  
16 volatile memory. Applications that generate volatile data on the IC module may  
17 place this volatile data in one or more data files within the volatile memory. The  
18 file system tracks the location of all data files as residing in either volatile memory  
19 or nonvolatile memory. When the same or a different application requests access  
20 to data, the file system determines whether a corresponding data file is located in  
21 volatile memory or nonvolatile memory and then locates the data file within the  
22 appropriate memory.

23 In this fashion, it is possible for multiple applications to share the same data  
24 in volatile memory by using the file system to access the volatile data files. By  
25 extending the file system to cover data in volatile memory, the volatile data files

1 are protected by the access control mechanisms of the file system so that only  
2 authorized applications can perform file operations on the volatile data files.

3 Generally, an IC module includes some form of processing capabilities, as  
4 well as limited volatile and nonvolatile memory. IC modules typically do not have  
5 their own power supply, but instead rely on an external power source, such as  
6 power provided by a card reader. In this manner, processing tasks are  
7 accomplished only when power is applied to the IC module. In other  
8 implementations, the IC modules may be implemented as USB Keys, built-into a  
9 motherboard of a computer or even inside of a larger microprocessor or ASIC  
10 (Application Specific Integrated Circuit). For discussion purposes, the IC module  
11 is described in the context of a smart card.

12 Fig. 1 shows an IC module 100 embodied as a smart card. The IC module  
13 100 has a reader interface 102 for interfacing with a card reader, a CPU or  
14 processor 104, a volatile rewritable memory in the form of RAM (Random Access  
15 Memory) 106, a ROM (Read Only Memory) 108, and an persistent read/write  
16 memory such as EEPROM (Electrically Erasable Programmable ROM) 110. A  
17 multi-bit bus (not shown) connects the components.

18 One or more applications 112 and an operating system 114 are stored in  
19 ROM 108. Some applications as well as parts of the operating system can reside  
20 in EEPROM as well. When the smart card is coupled to a card reader and receives  
21 power, the application(s) 112 and operating system 114 are executed on the  
22 processor 102. The operating system 114 exposes a set of application program  
23 interfaces (APIs) that enable resident applications 112 to perform tasks and  
24 manipulate data on the smart card. In addition, one or more nonresident  
25 applications 116, which execute external to the smart card (e.g., programs on

1 kiosks, point-of-purchase machines, etc.), may also place function calls with the  
2 operating system 114 to perform tasks or manipulate data on the smart card.  
3 Examples of such tasks include access security, cryptographic functions (e.g.,  
4 encryption, decryption, signing, and verification), file management, commerce,  
5 and so forth. One suitable operating system is the "Windows for Smart Card"  
6 operating system from Microsoft Corporation.

7 The operating system 114 includes a file system 118 that manages files  
8 stored on the smart card. Typically, smart card files are stored in nonvolatile  
9 memory, such as nonvolatile data files 120 in EEPROM 110. However, with this  
10 system, data in volatile memory may also be stored in special files, as represented  
11 by volatile data files 122 in RAM 106. The volatile files 122 make it possible for  
12 multiple resident applications 112, as well as nonresident applications 116 that are  
13 downloaded for a particular session, to share the same data in volatile memory 106  
14 (assuming the applications are authorized).

15 The file system 118 manages both the nonvolatile files 120 in EEPROM  
16 110 as well as the volatile files 122 in RAM 106. Reading and writing of data in  
17 both the nonvolatile memory 110 and the volatile memory 106 are accomplished  
18 by accessing the file(s) using the same APIs. This encapsulates the memory  
19 management layer and makes it easier for programmers to manipulate data in  
20 volatile memory.

21 The file system 118 enforces two levels of security for access to the volatile  
22 files 122. As a first level, the volatile files 122 in RAM 106 are protected by  
23 access control mechanisms of the file system 118. For example, an access control  
24 list (ACL) can be associated with each file in the file system 118, including both  
25 nonvolatile files 120 and volatile files 122. The ACL controls access to the

1 associated file, thereby permitting only authorized applications to gain access to  
2 and perform file operations on the volatile files 122. In this manner, unauthorized  
3 applications are unable to gain access to the data in the volatile files 122 of RAM  
4 106.

5 As a second level of security, the volatile files 122 are destroyed when  
6 power to the smart card is terminated (e.g. removing the smart card from the card  
7 reader) or a special command is received. When power is restored, the operating  
8 system initializes the file system by deleting any data in the volatile memory (if  
9 any) and removing any reference to volatile files that may exist in nonvolatile  
10 memory.

### 11 File System For Volatile And Nonvolatile Files

12 Fig. 2 illustrates a logical representation of the file system 118 implemented  
13 on IC module 100. The file system 118 manages access to both volatile and  
14 nonvolatile files in various types and regions of memory. In this example, the file  
15 system 118 facilitates access to volatile files 122(1)-122(4) in RAM 106,  
16 nonvolatile read/write files 120(1)-120(4) in EEPROM 110, and nonvolatile read-  
17 only files 200(1)-200(2) in ROM 108.

18 A set of file system APIs 202 allows applications to directly manipulate  
19 files and directories. Using the APIs, applications are able to create directories  
20 and files, as well as delete, close, read, and write to files. The same APIs are used  
21 for both volatile files 122 and nonvolatile files 120, 200. Thus, reading and  
22 writing of data in the volatile memory is accomplished by accessing the volatile  
23 data file(s) in RAM 106 using the same APIs used to access non-volatile  
24 read/write data files in EEPROM 110. One example set of basic APIs is provided  
25



below beneath the heading "Exemplary File System APIs". Other APIs can involve setting and getting file attributes (such as access permissions or file length), directory creation and deletion, and file deletion.

The file system 118 includes an ACL (access control list) table 204 that performs the security function of determining which users and/or applications have access to which files. The ACL table 204 might contain fields identifying what permissions have been granted to a corresponding file, the access control list that applies to the file, and the like. Access control lists and ACL tables are known in the art and are not described in detail.

The file system 118 also includes a memory region directory 206 that tracks the memory regions in which various files are located. The memory region directory 206 is illustrated in EEPROM 110 because the contents can be modified and updated as files are created, moved, or deleted. It also needs to be persistent over successive sessions, i.e. after the power is removed from the card. The memory region directory 206 includes a file name field to identify various files stored on the IC module and a memory identifier field to specify the corresponding memory region. In this example, file 1 is located in RAM 106, file 2 is located in ROM 108, and file 3 is located in EEPROM 110.

The file system 118 also includes a file location specifier to identify the physical location of the files within their relevant memory locations. In the illustrated implementation, the file system uses two different types of file location specifiers depending upon the memory region being accessed. For EEPROM 110 and RAM 106, the file system 118 employs a FAT file system, which garners its name for its method of organization—the file allocation table (FAT). Here, one file allocation table 210 resides in EEPROM 110 to locate nonvolatile read/write

1 files 120(1)-120(4) stored in EEPROM. A second file allocation table 212 is  
2 temporarily constructed in RAM 106 to locate volatile files 122(1)-122(4) that are  
3 stored in RAM 106. With FAT file systems, the data contained in a file need not  
4 be in contiguous areas of memory.

5 For ROM 108, the file system 118 employs a ROM file table 214 that  
6 identifies the locations of files 200(1) and 200(2) that are fixed in memory. The  
7 ROM file table stores the size of the file and its start address in the ROM. The file  
8 data is placed in contiguous areas of the ROM. Since files in ROM 108 are  
9 permanently fixed in the read only structure, the files cannot be deleted or resized.  
10 Thus, the advantages of a FAT file system are not applicable in ROM, resulting in  
11 the use of the ROM file table.

12 While the file system is described in the context of FAT and fixed location  
13 tables, it is noted that other types of file systems may be used to implement the  
14 features and methods described below, such as the ROM scheme.

15 When power is applied to the IC module, one or more applications may  
16 create and store volatile files 122(1)-122(4) in RAM 106. The file system 118  
17 constructs the RAM-based file allocation table 212 to track the volatile files. An  
18 entry for each volatile data file is then placed in the memory regions directory 206  
19 to indicate that the file is located in RAM.

20 One or more applications may access individual volatile files 122 using the  
21 APIs 202. For instance, suppose an application would like to open a volatile file  
22 named "File1". In an implementation using Windows-brand operating system, the  
23 application calls a function named "ScwCreateFile( )", which opens an existing  
24 file or creates a new file if one does not exist. In response to this function call, the  
25 file system 118 initially visits the ACL table 204 to determine whether the

1 application has sufficient authorization to access the requested file. This is  
2 represented by flow arrow 230 in Fig. 2.

3 Assuming the application is authorized, the file system 118 proceeds to the  
4 memory region directory 206 to determine which memory region holds the  
5 requested file (flow arrow 232). In this example, the directory 206 indicates that  
6 the file "File1" is located in RAM 106. The memory region directory 206 directs  
7 reference to the RAM-based file allocation table 212 (flow arrow 234), which  
8 provides physical location information to locate the volatile file 122(1) within  
9 RAM 106 (flow arrow 236).

10 The application may alternatively elect to open a nonvolatile file, such as  
11 "File2" or "File3". The application uses the same API function ScwCreateFile( )  
12 as it did for the volatile file. In the case of File2, the memory region directory 206  
13 determines that the file is located in ROM 108 and directs the file request to the  
14 ROM file table 214 (flow arrow 238). The ROM file table 214 specifies the  
15 physical location of the nonvolatile file 200(1) in ROM 108 (flow arrow 240). For  
16 File3, the memory region directory 206 directs the request to the EEPROM-based  
17 file allocation table 210 (flow arrow 242), which indicates the physical location of  
18 the nonvolatile file 120(1) in EEPROM 110 (flow arrow 244).

19 In this manner, the file system APIs 202 may be used to access files,  
20 regardless of where those files are located. That is, both volatile files and  
21 nonvolatile files may be accessed using the same APIs. Once the files are located,  
22 a handle to the file may be passed in to point to the exact location in memory,  
23 whether that location is in RAM, EEPROM, or ROM. The handles are obtained  
24 by creating a new file or opening an existing one, and may be passed in as a  
25 parameter of the function calls.

1 In a second implementation, the file system assumes that any file in a  
2 particular directory will reside in the volatile memory region. As an example, the  
3 file system understands that any file created in the "/r" directory is created in  
4 RAM. The same APIs are employed.

5 As a third implementation, the file system may be configured to identify the  
6 proper region of memory without using the memory region directory 206. In this  
7 implementation, the APIs are modified to include an additional parameter  
8 identifying the memory region in which the file is located. For example, the  
9 function ScwCreateFile( ) to create/open a file may be modified to include a  
10 parameter specifying whether the file is in RAM, ROM, or EEPROM, as follows:

11  
12 hFile = ScwCreateFile(fileName, aclFileName, CREATE\_IN\_RAM)

13  
14 Once created and accessed, a file handle may be used from that point on, as  
15 follows:

16  
17 ScwGetFileLength(hFile)

18  
19 With this third implementation, the application specifies the memory region  
20 where the file is to be created at the time of creation. This is not the case in the  
21 first implementation because the memory region directory tracks the memory  
22 region for the application.

23 With this architecture, volatile data kept in volatile memory is no longer  
24 bound to a single application, but can be accessed by multiple applications. Thus,  
25 one application can produce data and store that data in a volatile file in RAM 106.

1 A second application with appropriate authorization may subsequently access the  
2 same file and use or modify the data.

3 By extending the file system over the volatile memory space, the volatile  
4 files are protected by the same access control mechanisms used to protect  
5 nonvolatile files. Each access request is evaluated using the access control  
6 mechanisms, such as ACL lists, to determine whether the user/application making  
7 the request is authorized to gain access to the data.

8 When a session is finished and power is removed from the IC module, all  
9 volatile files 122 and the RAM-based FAT table 212 are immediately destroyed.  
10 Thus, this information will not be available upon restoration of power sensitive  
11 data is not left in tact within volatile memory, preventing any improper use of the  
12 data by a rogue application when power is restored. As a further measure of  
13 protection, when power is restored, the operating system 114 initializes the file  
14 system by deleting any data left in the volatile memory and removing any  
15 references in the file system that pertain to volatile files, such as entries in the  
16 memory region directory 206.

### 18 Operation

19 AI INS Fig. 3 illustrates a process 300 for operating the file system 118 of the IC  
20 module 100. At block 302, the file system 118 is initializes portions of the file  
21 system relating to the volatile memory. The file system 118 deletes any data in the  
22 volatile memory (e.g., RAM 106) and removes any entries in the memory region  
23 directory 206 pertaining to volatile files.

24 At block 304, in response to a request from an authorized application to  
25 create or open a file, the file system 118 creates or opens a file and obtains a

1 handle to that file. One exemplary process of opening a file is described below in  
2 more detail below with respect to Fig. 4. If a new file is created in RAM 106, a  
3 FAT table 212 is created in RAM 106 to track the location of the volatile file and  
4 an entry is added to the memory region directory 206 in EEPROM 110.

5 At block 306, the application is then free to read/write data to the volatile  
6 file. When the application is finished with the volatile file, it may request that the  
7 file be closed (block 308). In response, the file system 118 deletes the file from  
8 the RAM 106 and removes entries pertaining to the file from the memory region  
9 directory 206.

10 Fig. 4 shows a process 400 for creating/opening a file on the IC module  
11 100. At block 402, the file system receives a request from an application to open a  
12 file. At block 404, the file system 118 evaluates whether the application has  
13 appropriate authorization to access the file. As part of this validation process, the  
14 file system 118 checks the ACL table 204 for permissions and whether the  
15 application is approved by a controlling access control list.

16 If the application is not authorized (i.e., the "no" branch from block 406),  
17 access to the requested file is denied (block 408). On the other hand, if the  
18 application is authorized (i.e., the "yes" branch from block 406), the file system  
19 118 determines the memory region of the file using the memory region directory  
20 206 (block 410). For the illustrated implementation of Fig. 2, decision block 412  
21 presents three possible memory regions: RAM 106, ROM 108, or EEPROM 110.

22 If the file is located in RAM 106, it is a volatile file and the file system 118  
23 uses the RAM-based file allocation table 212 to locate the volatile file (block 414).  
24 If the file is located in EEPROM 110, it is a nonvolatile read/write file and the file  
25 system 118 uses the EEPROM-based file allocation table 210 to locate the

1 nonvolatile read/write file (block 416). If the file is located in ROM 108, it is a  
2 nonvolatile read-only file and the file system 118 uses the ROM file table 214 to  
3 locate the nonvolatile read-only file (block 418).

4 Once the file is located, the file system returns a file handle for use by the  
5 requesting application (block 420). At this point, the application may perform  
6 operations on the file, such as reading and writing (except in the case of a ROM-  
7 based read-only file).

### 8 9 Exemplary File System APIs

10 The file system APIs allow a smart card application to directly manipulate  
11 files and directories, including the ability to create directories and files as well as  
12 delete, close, read, and write to files. According to one implementation, the file  
13 system architecture extends the APIs to volatile files stored in volatile memory.  
14 Table 1 lists a set of exemplary file system APIs that are supported by the  
15 Windows® for Smart Card operating system from Microsoft. These APIs are used  
16 for both volatile data files stored in volatile memory and nonvolatile data files  
17 stored in nonvolatile memory.

18  
19 **Table 1**

20 <u>Functions</u>	21 <u>What it does...</u>
22 ScwCreateFile	Opens an existing file or creates a new file if one does not exist.
23 ScwCreateDir	Creates a directory with the given access control list (ACL) file.

ScwDeleteFile	Deletes the indicated file, which could be an access control list (ACL) file or a directory.
ScwCloseFile	Closes the file identified by the file handle.
ScwReadFile	Reads the indicated number of bytes at the current offset.
ScwWriteFile	Writes the indicated number of bytes at the current offset.
ScwGetFileLength	Returns the file size.
ScwGetFileAttributes	Returns the file attribute value associated with the given file.
ScwSetFileAttributes	Sets the file attribute value associated with a given file.
ScwSetFilePointer	Changes the current offset of the file to the indicated offset.
ScwSetFileACL	Sets the access control list (ACL) file for the given file.
ScwGetFileACLHandle	Opens the associated access control list (ACL) file for a given file.

## **Conclusion**

Although the invention has been described in language specific to structural features and/or methodological steps, it is to be understood that the invention defined in the appended claims is not necessarily limited to the specific features or steps described. Rather, the specific features and steps are disclosed as preferred forms of implementing the claimed invention.